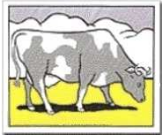


An Evaluation of Server Consolidation Workloads for Multi-Core Designs

Natalie Enright Jerger,
Dana Vantrease, Mikko H. Lipasti

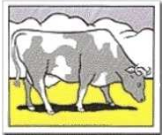
University of Wisconsin - Madison

IEEE International Symposium on Workload Characterization
September 27, 2007



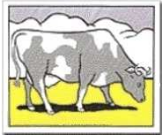
Motivation

- n Multi-core chips are here
- n Many-core chips (10-100s of cores) are on the horizon
 - n We need applications to evaluate future systems now
 - n Programming parallel applications is not getting any easier
- n Many-core systems can serve as consolidation platforms for multiple discrete applications



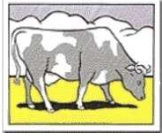
Server Farms

- n Consequences of growing server farms
 - n Higher management costs/overheads
 - n Higher floor space and electricity consumption
 - n Less reliable
 - n Individual servers may be underutilized
- n Virtualization facilitates consolidation of several physical servers onto a single high-end system
 - n Reduces management costs/overheads
 - n Increases overall utilization



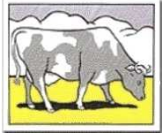
Server Consolidation Overview

- n Server workload characteristics well suited to multi-core architectures
 - n Multithreaded
 - n Take advantage of abundance of on-chip cores
 - n Communication intensive
 - n Low latency cache-to-cache transfers
 - n On-chip sharing
- n Potential interference between applications
 - n Interconnect bandwidth
 - n Memory controllers
 - n Cache resources



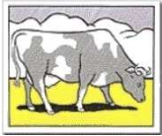
Opportunities with Server Consolidation Workloads

- n Explore the relevance of server consolidation workloads to a variety of research communities
- n Identify the interactions within server consolidation workloads
 - n Open up new avenues of research
 - n We focus on cache sharing within and between applications



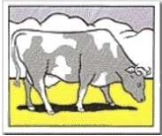
Outline

- n Motivation
- n Workloads
- n Cache Configurations
- n Evaluation Methodology
- n Results
- n Conclusions



Workloads

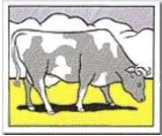
Workloads	Description	Setup	Execution
SPECjbb	Order processing application for wholesaler. Emphasizes the middle-tier business logic and performance of Java-based middleware.	3-tier client-server with 6 warehouses	6400 requests with 15 seconds of warm-up time
SPECweb	World-wide web server	3-tier, Zeus Web Server 3.3.7	300 HTTP requests
TPC-H	Decision Support System	IBM DB2 v6.1	Query #12 (shipping modes and order priority) on 512 MB database with 1GB of memory
TPC-W	Web commerce modeling online bookstore	IBM DB2 v6.1	Browsing mix for 25 web transactions



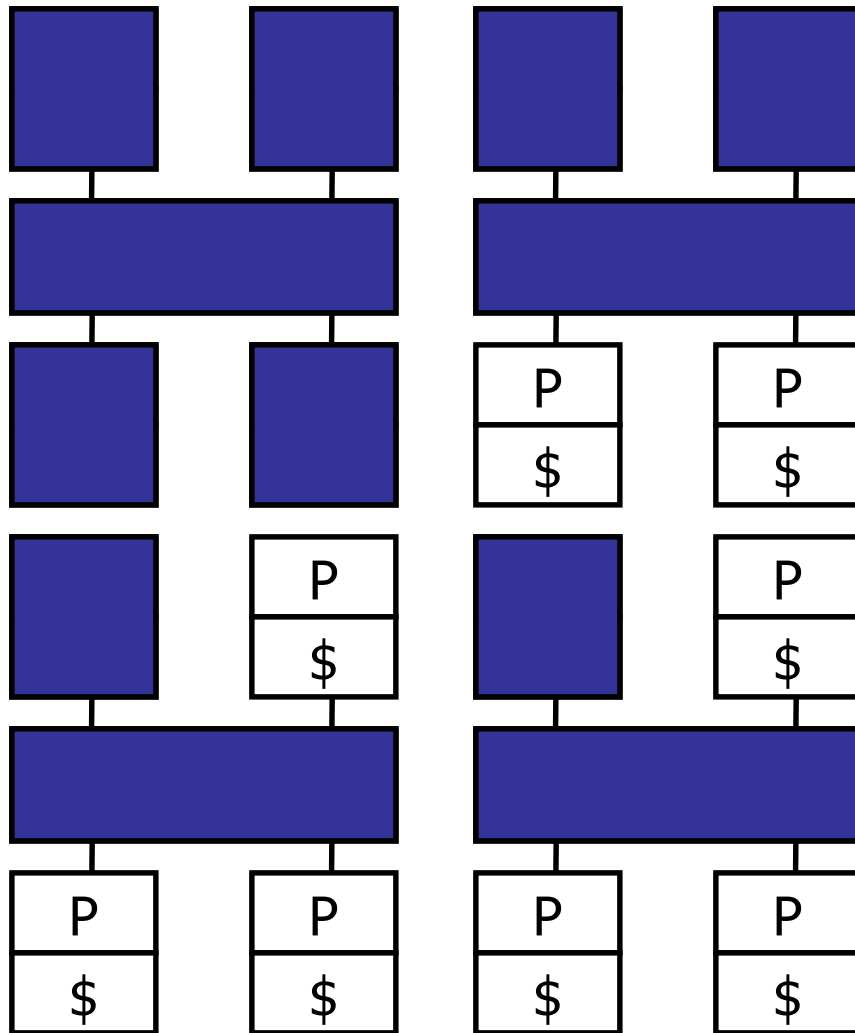
Workload Characteristics

	Percent of accesses resulting in \$-to-\$ transfer			Memory Footprint
	Clean	Dirty	Clean+Dirty	
SPECjbb	13%	2%	15%	Large
SPECweb	49%	3%	52%	Medium
TPC-H	30%	39%	69%	Small
TPC-W	34%	3%	37%	Large

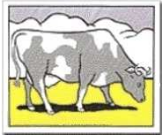
- n Many different combinations possible



Target CMP

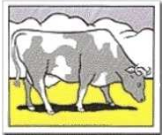


- n Affinity
 - n Maximizes sharing
- n Round Robin/Interleaved
 - n Maximizes available cache
- n Round Robin – Affinity Hybrid
- n Random
 - n Result of load balanced scheduling



Cache Configurations

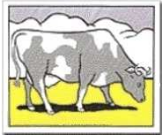
- n Vary degree of sharing
 - n Private
 - n Partially shared
 - n Shared
- n Vary assignments of threads to cores and caches
 - n Affinity
 - n Round Robin
 - n RR-Affinity Hybrid
 - n Random



Simulation Methodology (1)

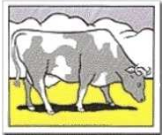
- n PHARMSim
 - n Full-system simulator
 - n Runs AIX OS
- n Methodology
 - n Each application isolated in its own Virtual Machine
 - n Runs its own copy of OS
 - n Statically assign a portion of global memory – private address space





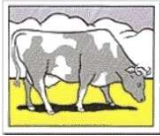
Simulation Methodology (2)

- n Each VM:
 - n Loads 4-processor checkpoint
 - n Snapshot of the workload
 - n System already booted and warmed
 - n Alleviates overhead of booting the OS
 - n Runs both user and operating system code
 - n Fixed number of transactions
 - n Ensures that the same set of transactions are run



Simulation Methodology (3)

- n At startup the system:
 - n Loads four VMs
 - n Each VM runs a pre-made 4-processor checkpoint
 - n Allows a mix and match of VMs for a variety of workload combinations
 - n Assigns a VM to a subset of cores
 - n All of a VM's threads run within their assigned domain
 - n Assignment is maintained throughout simulation

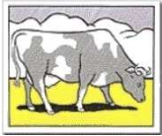


Workload Mixes

Heterogeneous Mixes

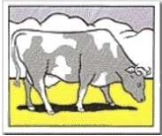
Mix 1	TPC-W (3) + TPC-H (1)
Mix 2	TPC-W (2) + TPC-H (2)
Mix 3	TPC-W (1) + TPC-H (3)
Mix 4	SPECjbb (3) + TPC-H (1)
Mix 5	SPECjbb (2) + TPC-H (2)
Mix 6	SPECjbb (1) + TPC-H (3)
Mix 7	SPECjbb (3) + TPC-W (1)
Mix 8	SPECjbb (2) + TPC-W (2)
Mix 9	SPECjbb (1) + TPC-W (3)

§ Homogeneous Mixes of TPC-W, TPC-H, SPECjbb, and SPECweb

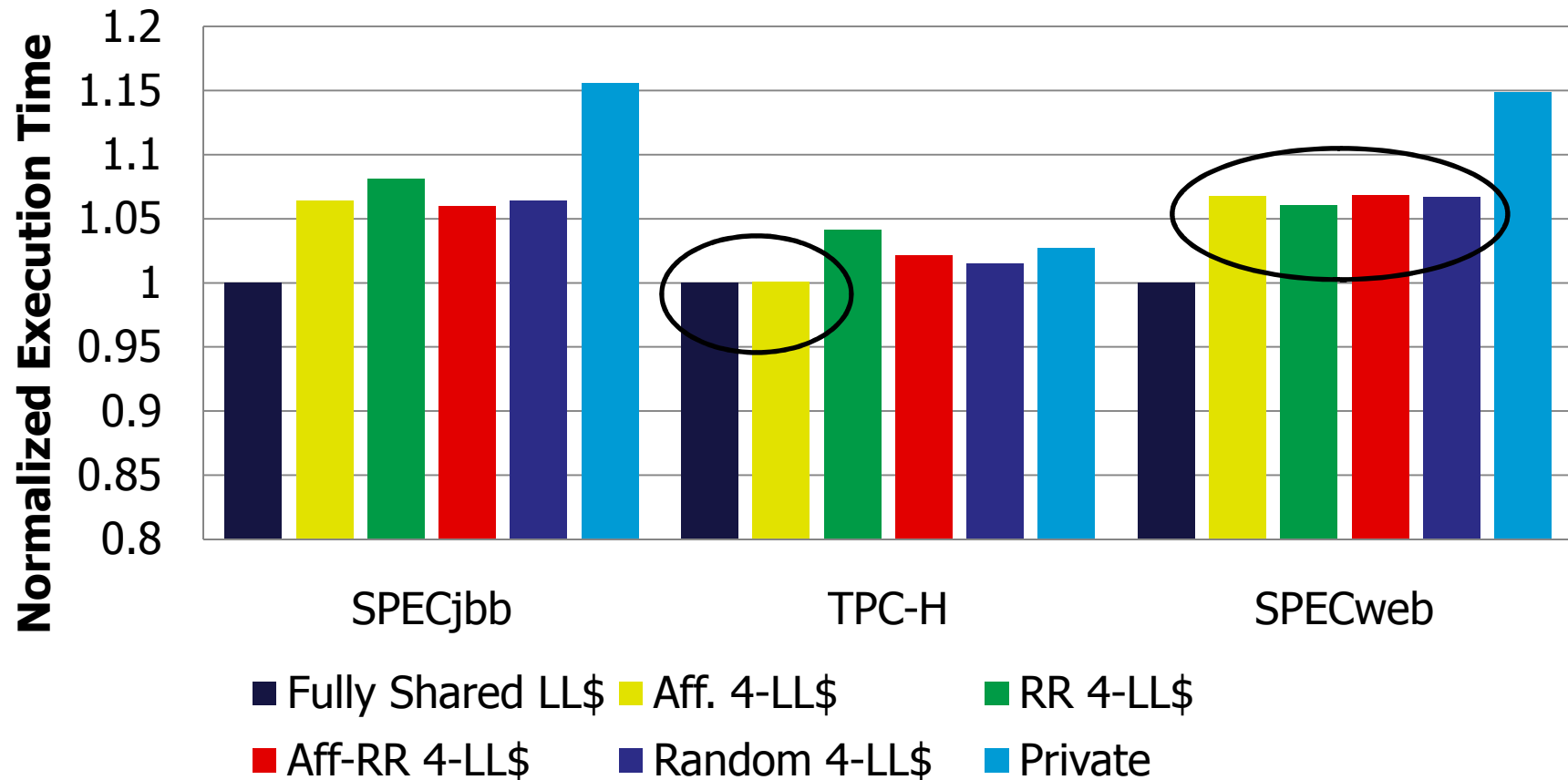


Simulation Parameters

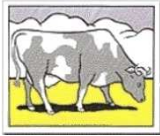
Machine Configuration	
Core	16 single-threaded, in-order
Interconnect	2-D Packet Switched Mesh
L0 Cache	Split I/D, 8 KB each
L1 Cache	Private 64 KB
L2 Cache	Shared, 16MB
Memory Latency	150 cycles
Cache Coherence	SGI Origin Directory Protocol
Thread to Core Assignment	Round Robin, Affinity, Round Robin-Affinity Hybrid, Random



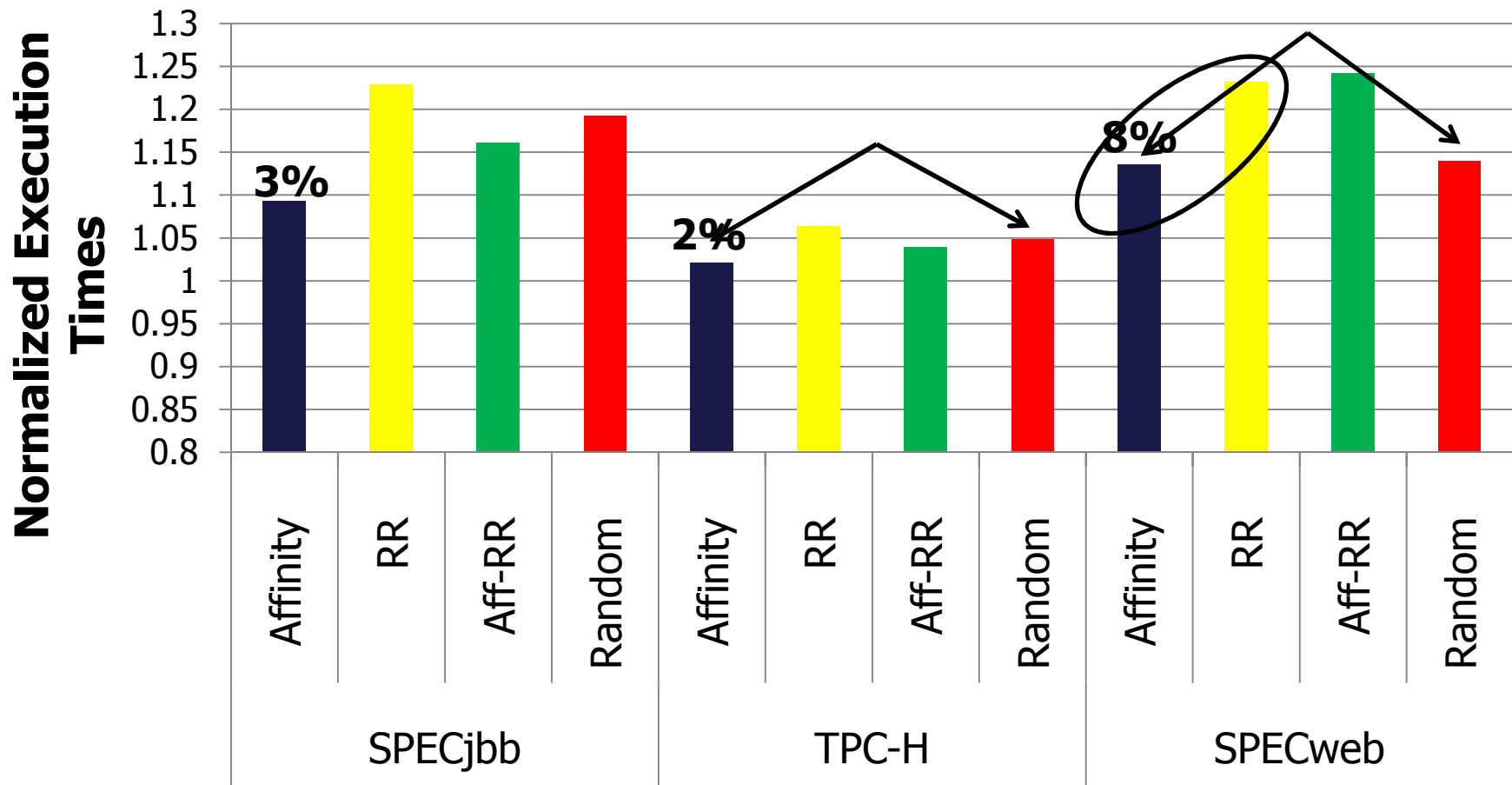
Isolated Workload Performance



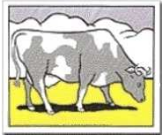
§ Affinity mapping *almost* always the best choice



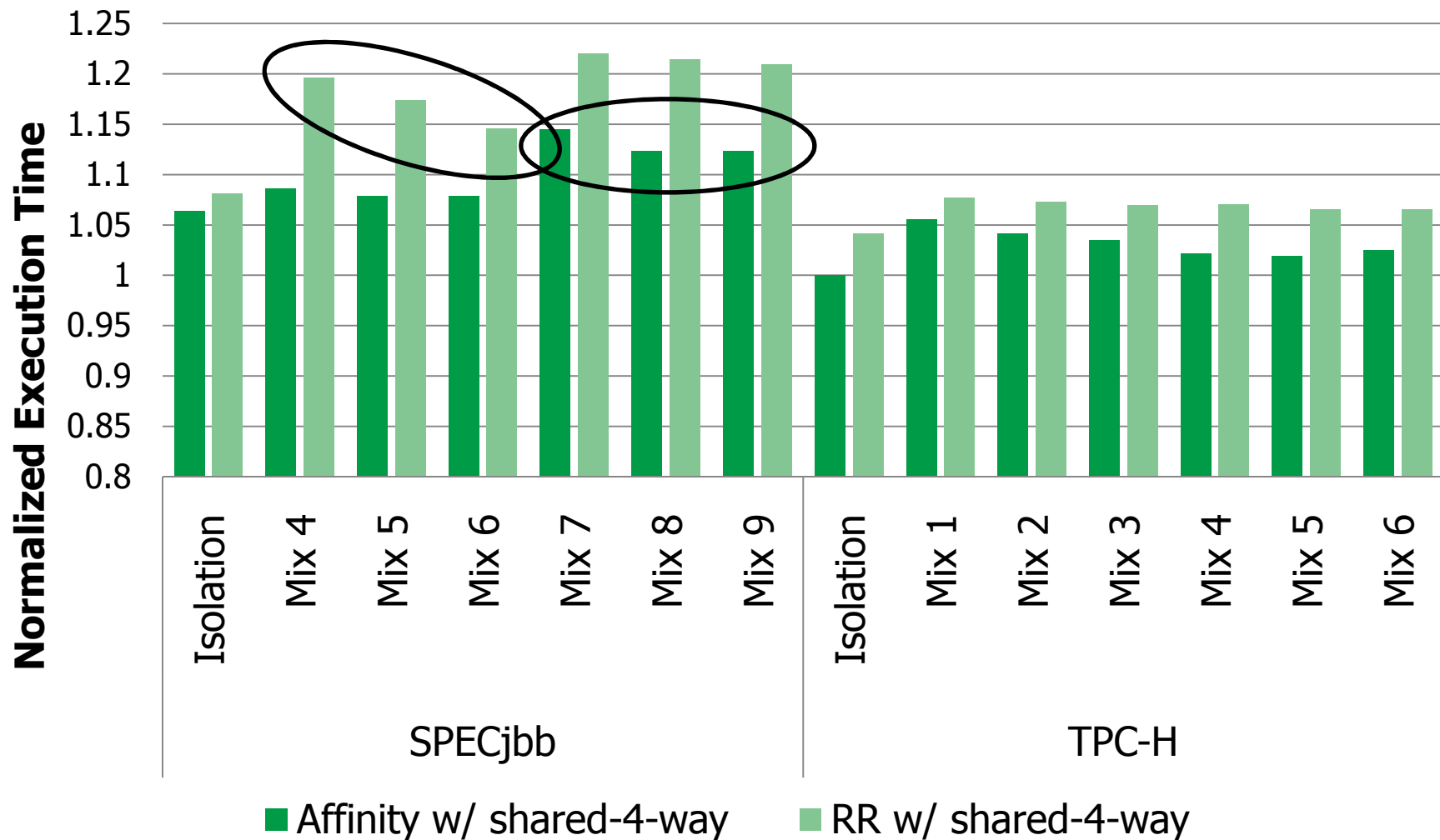
Homogeneous Mix Performance

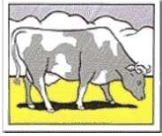


§ Performance differences between placement policies become more pronounced in consolidated environment



Heterogeneous Mix Performance





Implications and Related Work

- n Quality of Service

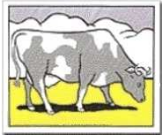
- n [Iyer, Sigmetrics 07], [Nesbit, MICRO-39]

- n Coherence Protocols

- n Virtual Hierarchies [Marty, ISCA 07]

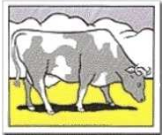
- n Methodology

- n SMT Methodology



Conclusions

- n Server consolidation workloads provide a framework on which to evaluate many-core architectures now
 - n Illustrate this by looking at degrees of cache sharing
- n Open up new and interesting avenues of research
 - n Architecture, Operating Systems, Virtual Machines



Questions?
